

# Hashing as a Dictionary Implementation

## Chapter 22



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# Contents

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# Objectives

- Describe relative efficiencies of various collision resolution techniques
- Describe hash table's load factor
- Describe rehashing and why necessary
- Use hashing to implement ADT dictionary

# Efficiency of Hashing

- Observations
  - Successful retrieval or removal has same efficiency as successful search
  - Unsuccessful retrieval or removal has same efficiency as unsuccessful search
  - A successful addition has same efficiency as unsuccessful search
  - An unsuccessful addition has same efficiency as successful search

# Efficiency of Hashing

- Load factor  $\lambda = \frac{\text{Number of entries in the dictionary}}{\text{Number of locations in the hash table}}$
- Minimum load factor = 0
  - When dictionary is empty
  - It is never negative
- Maximum load factor
  - Depends on type of collision resolution used
  - Cannot exceed 1



# Open Addressing

- Average number of collisions for *linear* probing

- Unsuccessful search  $\frac{1}{2} \left\{ 1 + \frac{1}{(1-\lambda)^2} \right\}$

- Successful search  $\frac{1}{2} \left\{ 1 + \frac{1}{(1-\lambda)} \right\}$

$\lambda$	Unsuccessful Search	Successful Search
0.1	1.1	1.1
0.3	1.5	1.2
0.5	2.5	1.5
0.7	6.1	2.2
0.9	50.5	5.5

Figure 22-1 The average number of comparisons required by a search of the hash table for given values of the load factor  $\lambda$  when using linear probing



# Open Addressing

- Average number of collisions for *quadratic probing* or *double hashing*
  - Unsuccessful search  $\frac{1}{(1-\lambda)}$
  - Successful search  $\frac{1}{\lambda} \log \left( \frac{1}{1-\lambda} \right)$

$\lambda$	Unsuccessful Search	Successful Search
0.1	1.1	1.1
0.3	1.4	1.2
0.5	2.0	1.4
0.7	3.3	1.7
0.9	10.0	2.6

Figure 22-2 The average number of comparisons required by a search of the hash table for given values of the load factor  $\lambda$  when using either quadratic probing or double hashing

# Separate Chaining

- Load factor is  $\lambda = \frac{\text{Number of entries in the dictionary}}{\text{Number of chains}}$
- Average number of comparisons
  - Unsuccessful search  $\lambda$
  - Successful search  $1 + \frac{\lambda}{2}$

$\lambda$	Unsuccessful Search	Successful Search
0.1	0.1	1.1
0.3	0.3	1.2
0.5	0.5	1.3
0.7	0.7	1.4
0.9	0.9	1.5
1.1	1.1	1.6
1.3	1.3	1.7
1.5	1.5	1.8
1.7	1.7	1.9
1.9	1.9	2.0
2.0	2.0	2.0

Figure 22-3 The average number of comparisons required by a search of the hash table for given values of the load factor  $\lambda$  when using separate chaining

# Rehashing

- When load factor gets too high
  - Resize array to a prime number at least twice former size
  - Must rehash to different locations  $c \% n$ , based on new size,  $n$ , of array
- Note – this is more work than simply increasing the size of an array
  - Not a task to be done often

# Comparing Schemes

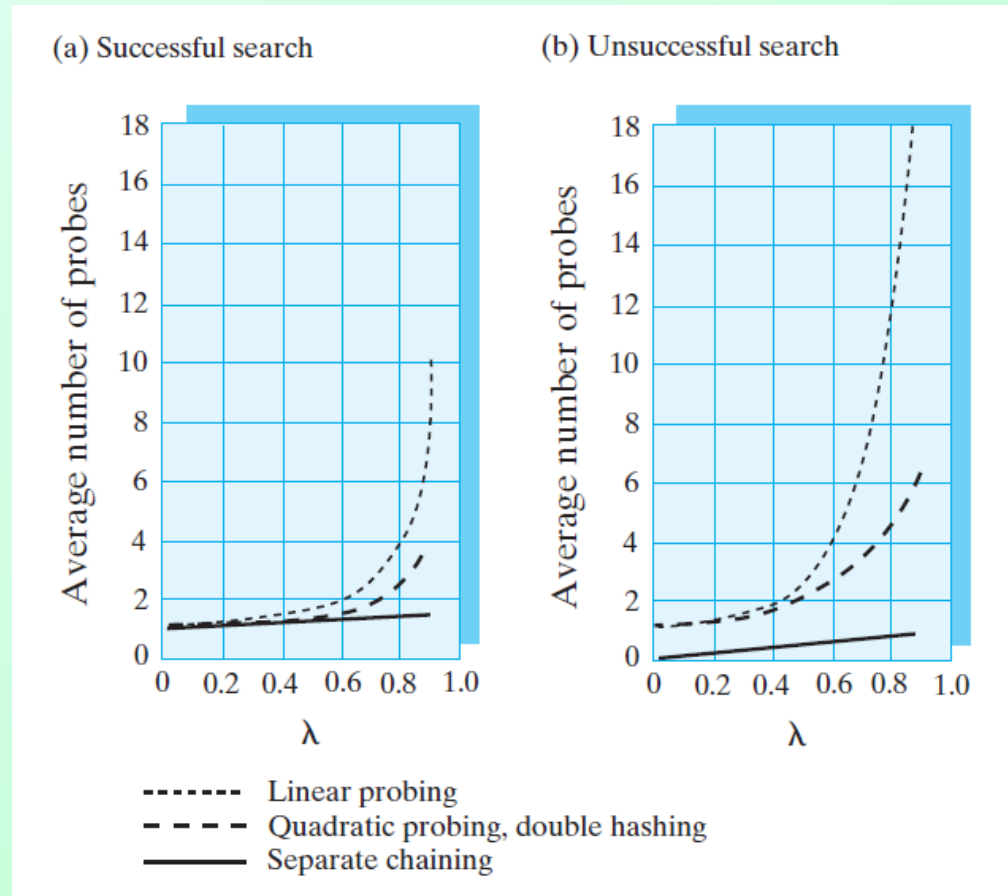


Figure 22-4 The average number of comparisons required by a search of the hash table versus the load factor  $\lambda$  for four collision resolution techniques when the search is (a) successful; (b) unsuccessful



# Dictionary Implementation That Uses Hashing

- We implement linear probing
  - Other open addressing strategies involve few changes
- Note source code of class **HashedDirectory**, [Listing 22-1](#)

Note: Code listing files  
must be in same folder  
as PowerPoint files  
for links to work

# Dictionary Implementation That Uses Hashing

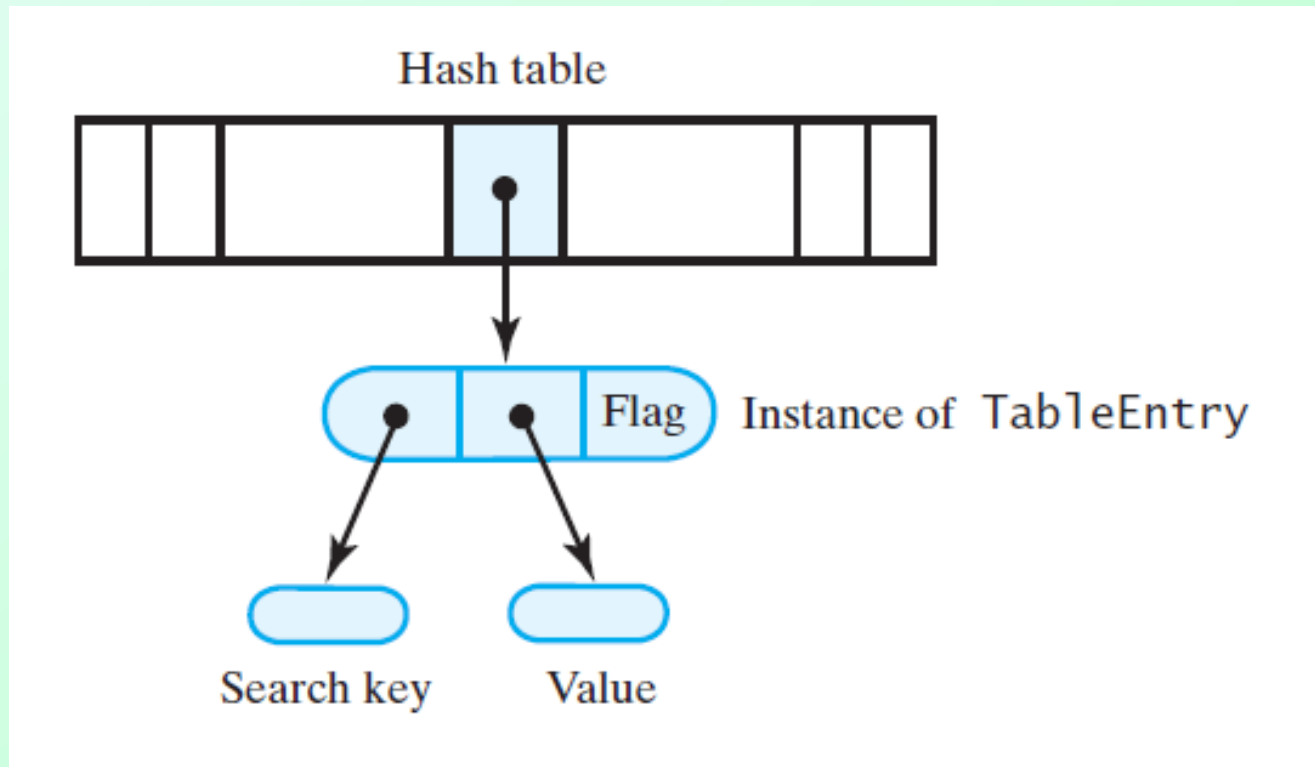


Figure 22-5 A hash table and one of its entry objects



Blue = current entry  
Light gray = removed entry  
Dark gray = null

FIGURE 22-6 A hash table containing dictionary entries, removed entries, and **null** values

# Java Class Library: The Class `HashMap`

- Standard package `java.util` contains the class `HashMap<K, V>`
- Table is a collection of buckets
- Constructors
  - `public HashMap()`
  - `public HashMap(int initialSize)`
  - `public HashMap(int initialSize, float maxLoadFactor)`
  - `public HashMap(Map<? extends K, ? extends V> table)`

# Java Class Library: The Class `HashMap`

- Design
  - Max  $\lambda = 0.75$
  - Avoid necessity of rehashing by setting

$$\text{Number of buckets} > \frac{\text{Max entries in dictionary}}{\lambda_{\max}}$$

# Java Class Library: The Class `HashSet`

- Implements the interface `java.util.Set` of Chapter 1
  - Uses an instance of the class `HashMap` to contain entries in a set
- Constructors
  - `public HashSet()`
  - `public HashSet(int initialCapacity)`
  - `public HashSet(int initialCapacity, float loadFactor)`



# End

## Chapter 22

