

# Sorted Lists

## Chapter 16



# Contents

- Specifications for the ADT Sorted List
  - Using the ADT Sorted List
- A Linked Implementation
  - The Method **add**
  - The Efficiency of the Linked Implementation
- An Implementation That Uses the ADT List
  - Efficiency Issues

# Objectives

- Use sorted list in a program
- Describe differences between ADT list and ADT sorted list
- Implement ADT sorted list by using chain of linked nodes
- Implement ADT sorted list by using operations of ADT list

# Sorted Lists

- We extend the capability of a list
  - Previous example used list to organize names in alphabetical order
- Consider need to keep list sorted in numerical or alphabetic order after list established
  - We add or remove an element
  - The ADT handles keeping elements in order

# Specifications for the ADT Sorted List

- Possible operations
  - For simplicity, duplicate entries allowed
  - Must determine where in list element is added
  - Can ask if list contains specified entry
  - Must be able to remove an entry

# Abstract Data Type: Sorted List

- Data
  - A collection of objects in sorted order and having same data type
  - Number of objects in collection
- Operations
  - `add(newEntry)`
  - `remove(anEntry)`
  - `getPosition(anEntry)`

# Abstract Data Type: Sorted List

- Operations used from ADT list (Ch. 12)
  - `getEntry(givenPosition)`
  - `contains(anEntry)`
  - `remove(givenPosition)`
  - `clear()`
  - `getLength()`
  - `isEmpty()`
  - `toArray()`
- Interface, [Listing 16-1](#)

Note: Code listing files must be in same folder as PowerPoint files for links to work

Question 1 Suppose that `wordList` is an unsorted list of words. Using the operations of the ADT list and the ADT sorted list, create a sorted list of these words.

Question 2 Assuming that the sorted list you created in the previous question is not empty, write Java statements that

a. Display the last entry in the sorted list.

b. Add the sorted list's first entry to the sorted list again.



Question 1 Suppose that `wordList` is an unsorted list of words. Using the operations of the ADT list and the ADT sorted list, create a sorted list of these words.

```
SortedListInterface<String> sortedWordList = new SortedList<String>();  
int numberOfWords = wordList.getLength();  
for (int position = 1; position <= numberOfWords; position++)  
    sortedWordList.add(wordList.getEntry(position));
```

Question 2 Assuming that the sorted list you created in the previous question is not empty, write Java statements that

a. Display the last entry in the sorted list.

```
int length = sortedWordList.getLength();  
String lastEntry = sortedWordList.getEntry(length);  
System.out.println(lastEntry);
```

b. Add the sorted list's first entry to the sorted list again.

```
sortedList.add(sortedList.getEntry(1));
```

# A Linked Implementation

- Linked implementation of the ADT sorted list, [Listing 16-2](#)
- Note different versions of method **add**
  - [Iterative](#) version
  - [Recursive](#) version

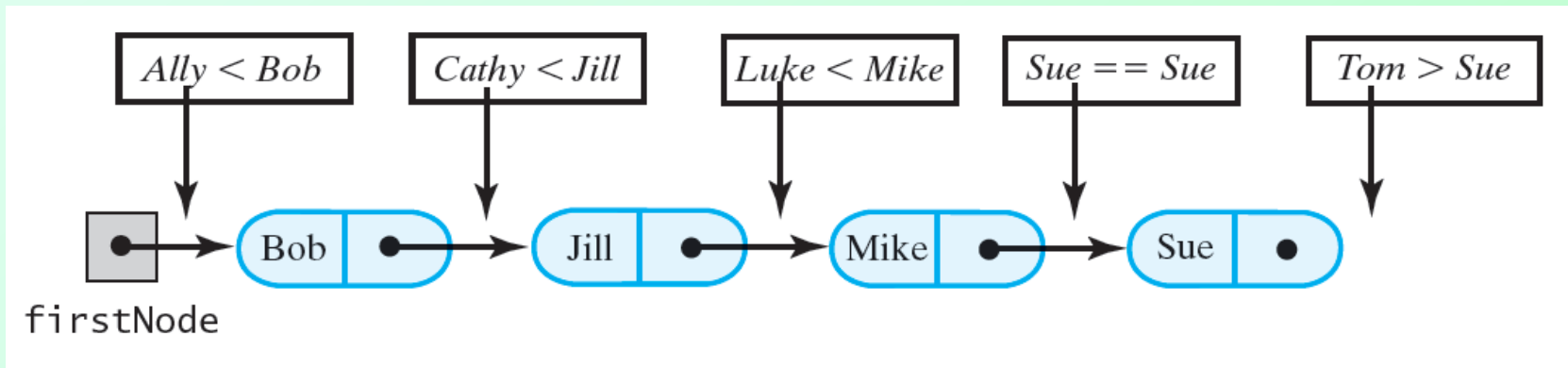


Figure 16-1 Places to insert names into a sorted chain of linked nodes

Question 3 In the while statement of the method `getNodeBefore` , how important is the order of the two boolean expressions that the operator `&&` joins? Explain.

```
while ((currentNode != null) && (anEntry.compareTo (currentNode.getData ())) > 0))
```

Question 4 What does `getNodeBefore` return if the sorted list is empty? How can you use this fact to simplify the implementation of the method `add` given in Segment 16.10?

Question 5 Suppose that you use the previous method `add` to add an entry to a sorted list. If the entry is already in the list, where in the list will `add` insert it? Before the first occurrence of the entry, after the first occurrence of the entry, after the last occurrence of the entry, or somewhere else?

Question 6 What would be the answer to the previous question if you changed `>` to `>=` in the while statement of the method `getNodeBefore` ?

Question 3 In the while statement of the method `getNodeBefore`, how important is the order of the two boolean expressions that the operator `&&` joins? Explain.

```
while ((currentNode != null) && (anEntry.compareTo (currentNode.getData ()) > 0))
```

The order is critical. When `currentNode` is null, `currentNode != null` is false. Thus, the entire expression in the while statement is false without executing the call `currentNode.getData()`. If the latter call were to execute first when `currentNode` was null, an exception would occur. Thus, the while statement should remain as written.

Question 4 What does `getNodeBefore` return if the sorted list is empty? How can you use this fact to simplify the implementation of the method `add` given in Segment 16.10?

When the sorted list is empty, `getNodeBefore` returns null. Thus, in the definition of `add`, you can omit the call to `isEmpty` in the if statement.

Question 5 Suppose that you use the previous method `add` to add an entry to a sorted list. If the entry is already in the list, where in the list will `add` insert it? Before the first occurrence of the entry, after the first occurrence of the entry, after the last occurrence of the entry, or somewhere else?

Before the first occurrence of the entry.

Question 6 What would be the answer to the previous question if you changed `>` to `>=` in the while statement of the method `getNodeBefore` ?

After the last occurrence of the entry.

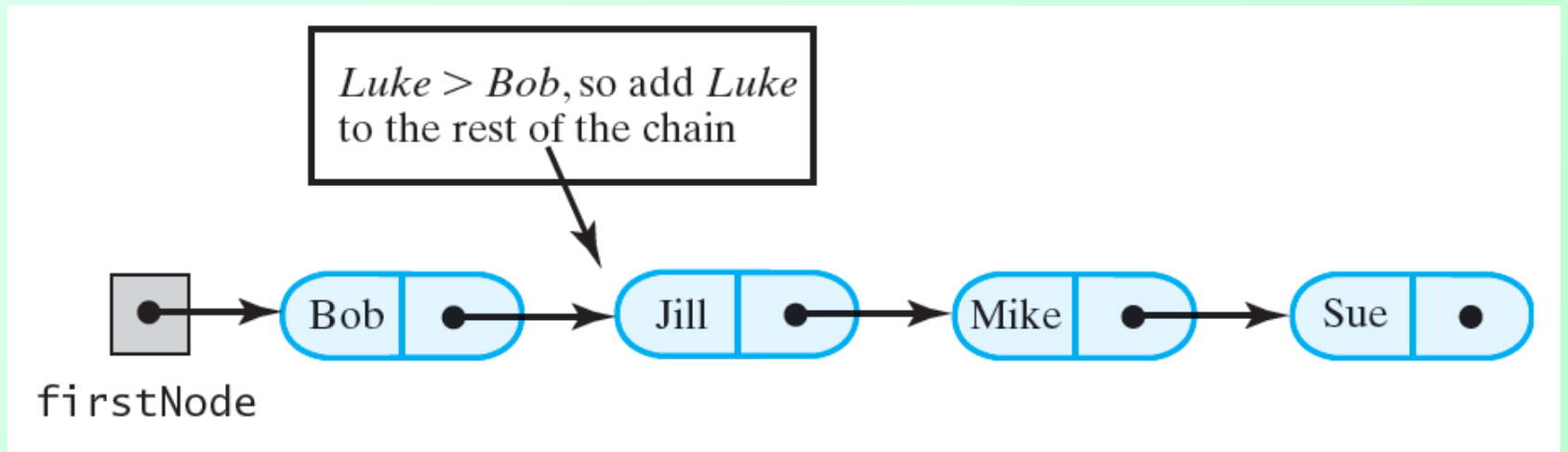


Figure 16-2 Recursively adding *Luke* to a sorted chain of names

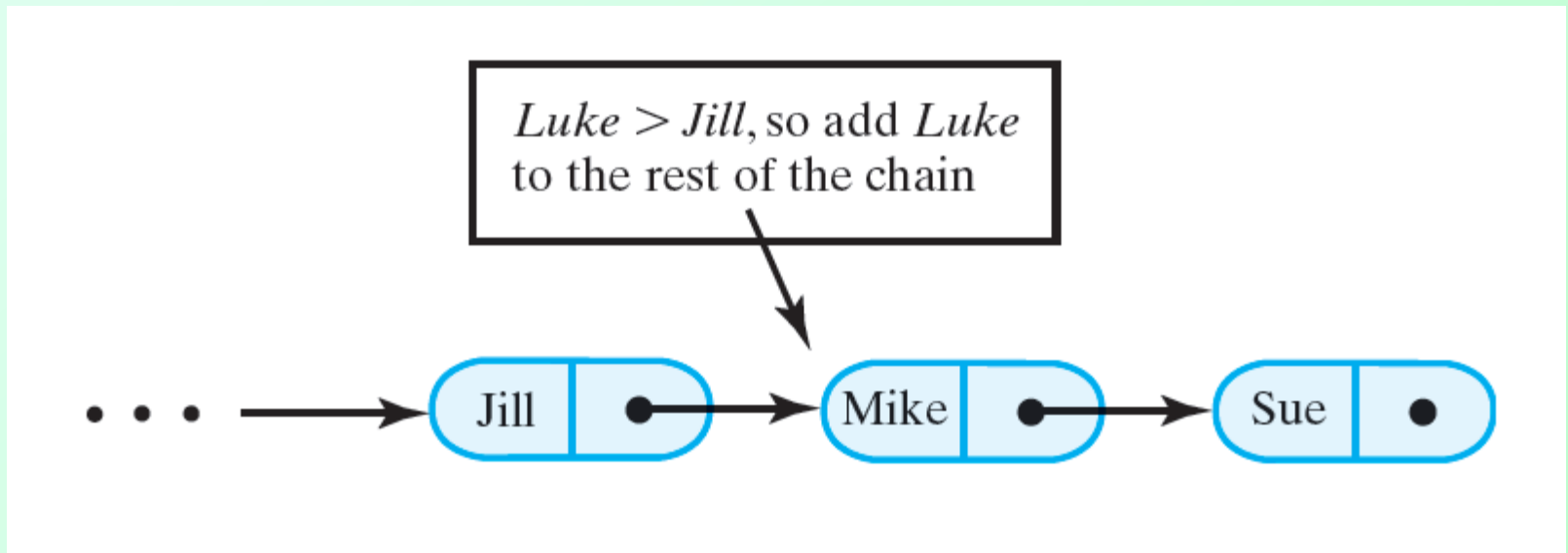


Figure 16-2 Recursively adding *Luke* to a sorted chain of names

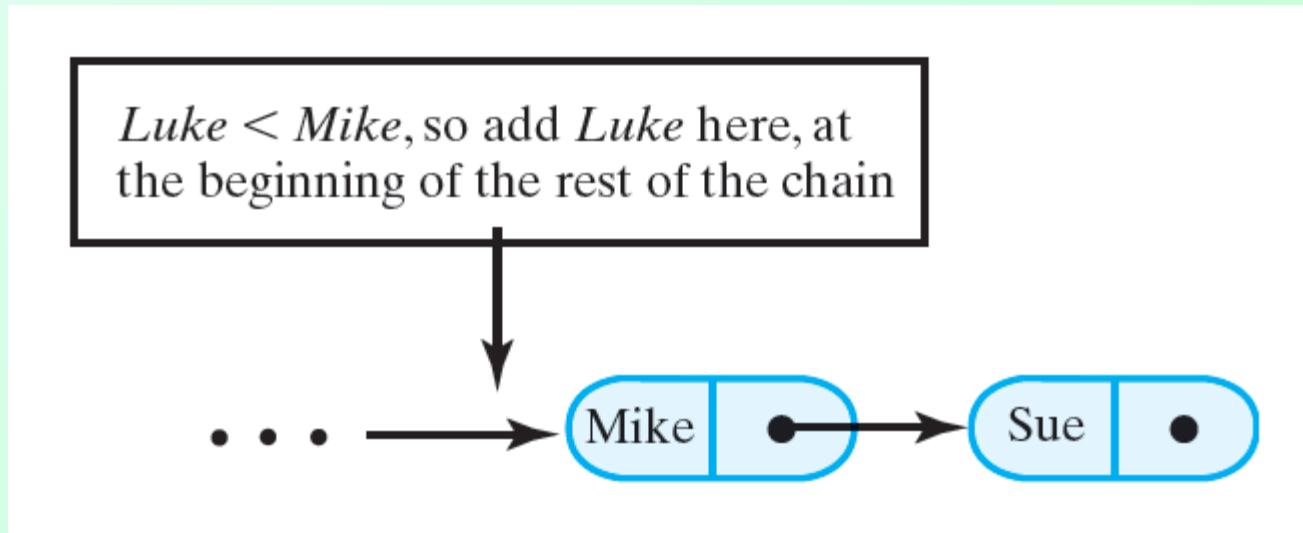


Figure 16-2 Recursively adding *Luke* to a sorted chain of names



(a) The list before any additions



(b) As `add("Ally", firstNode)` begins execution

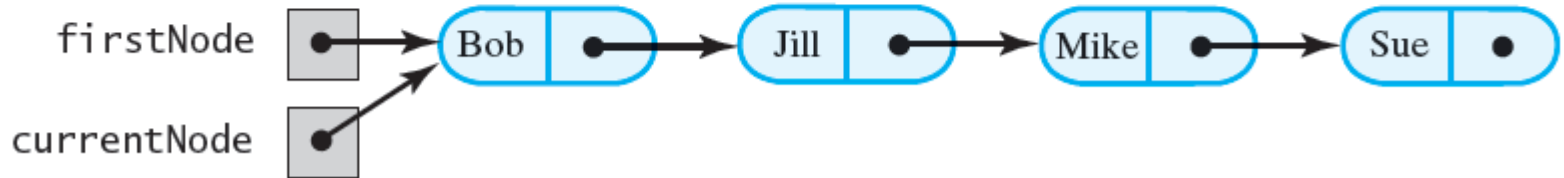
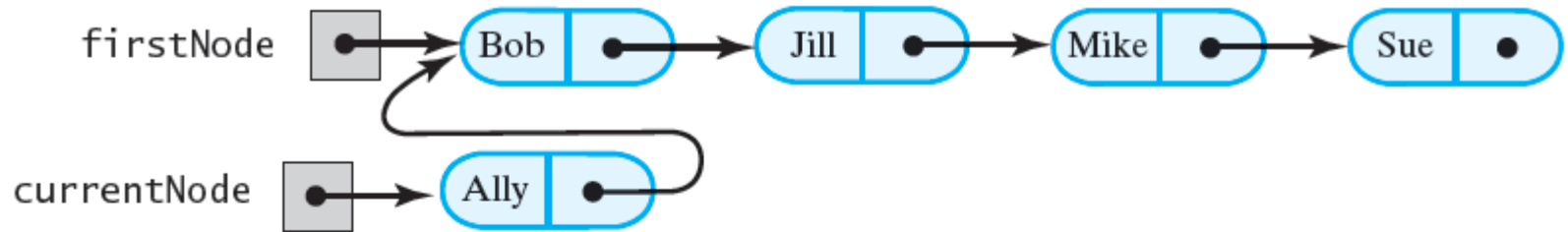


Figure 16-3 Recursively adding a node at the beginning of a chain

(c) After a new node is created (the base case)



The private method returns the reference that is in `currentNode`

(d) After the public `add` assigns the returned reference to `firstNode`

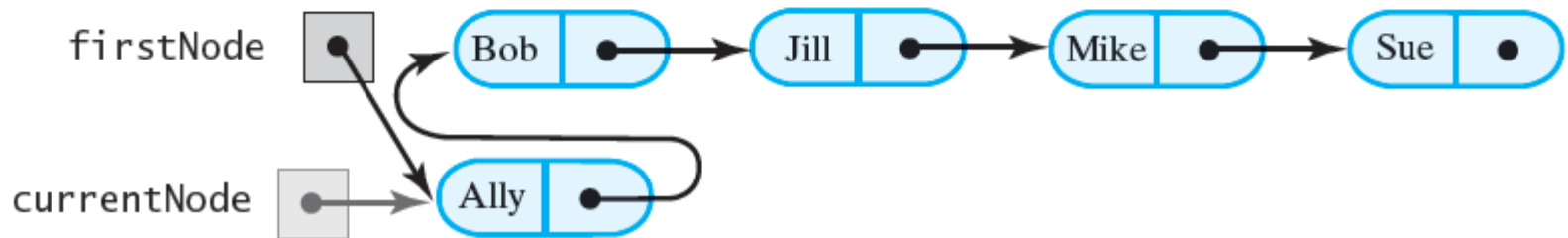
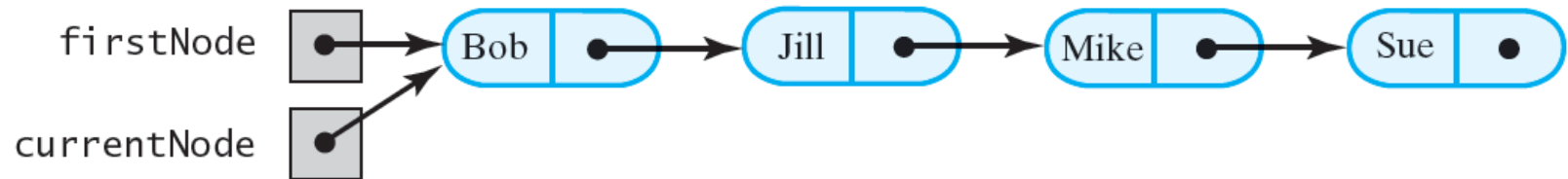


Figure 16-3 Recursively adding a node at the beginning of a chain

(a) As `add("Luke", firstNode)` begins execution



(b) As the recursive call `add("Luke", currentNode.getNextNode())` begins execution

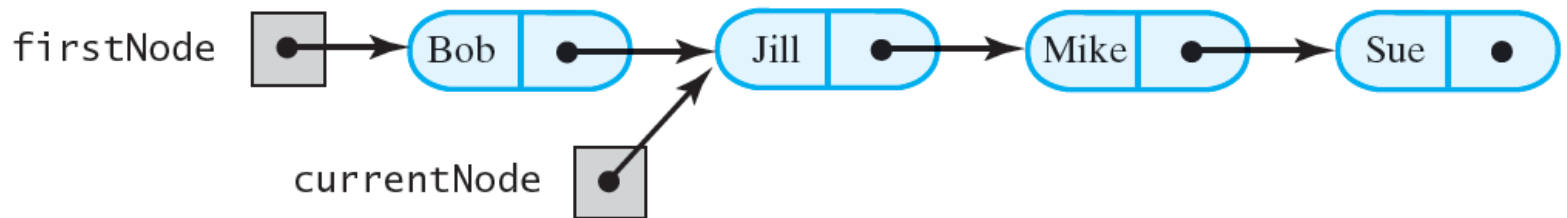
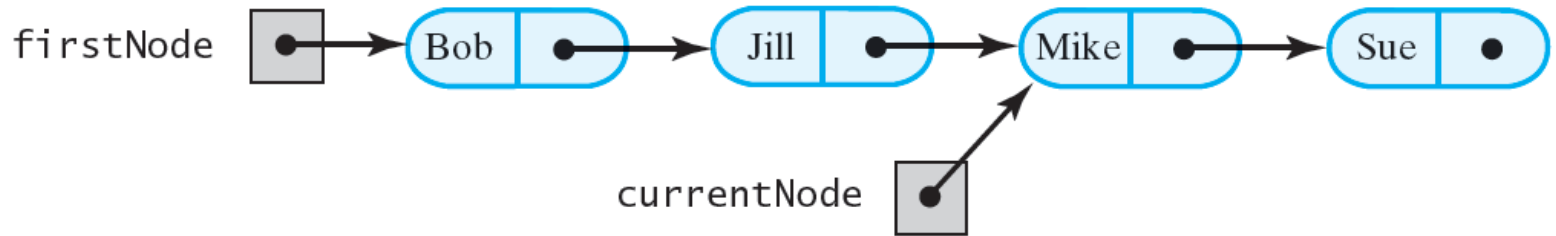
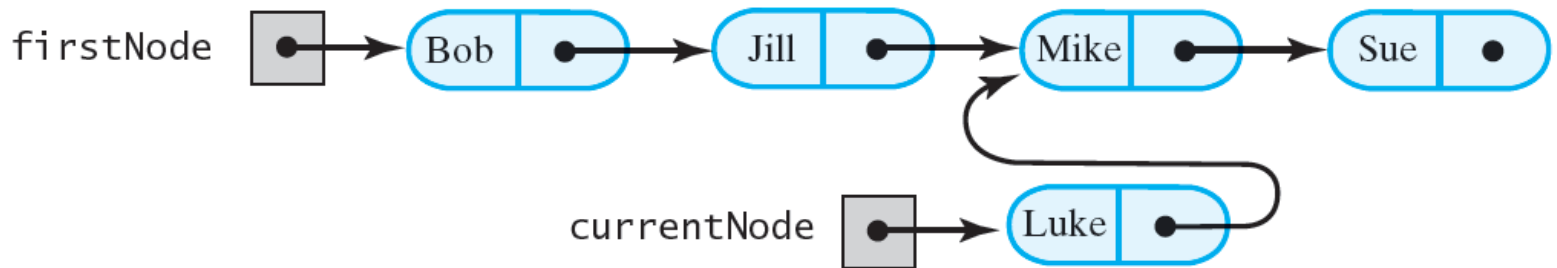


Figure 16-4 Recursively adding a node between existing nodes in a chain

(c) As the recursive call `add("Luke", currentNode.getNextNode())` begins execution



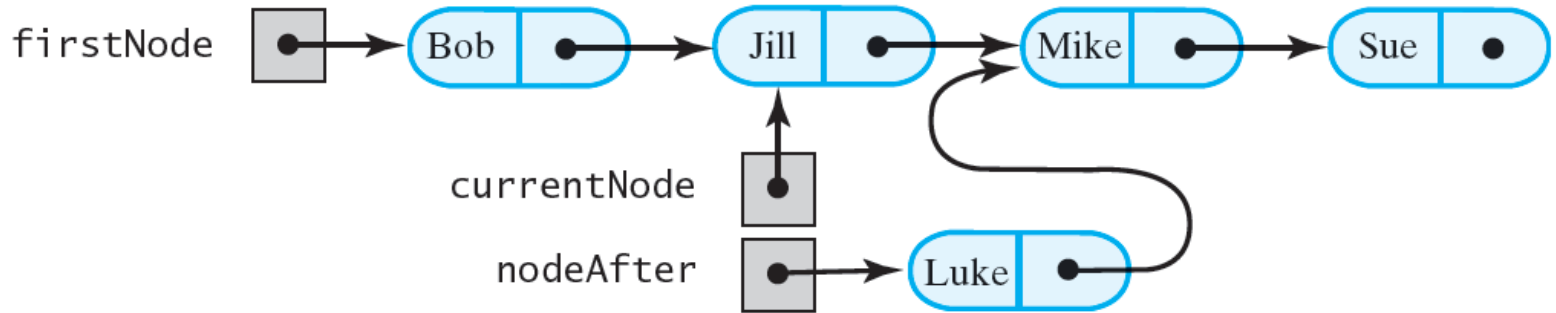
(d) After a new node is created (the base case)



The private method returns the reference that is in `currentNode`

Figure 16-4 Recursively adding a node between existing nodes in a chain

(e) After the returned reference is assigned to `nodeAfter`



(f) After `currentNode.setNextNode(nodeAfter)` executes

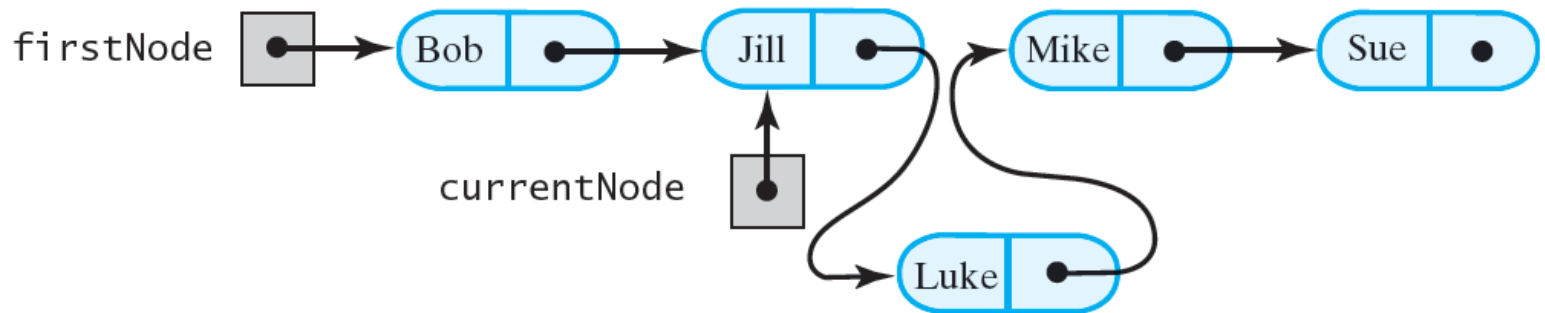


Figure 16-4 Recursively adding a node between existing nodes in a chain

Question 8 The linked implementation of the ADT sorted list, as given in this chapter, does not maintain a tail reference. Why is a tail reference more significant for a linked implementation of the ADT list than it is for a sorted list?

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The method `add(newEntry)` for a list adds a new entry at the end of the list. A tail reference makes this method  $O(1)$  instead of  $O(n)$ . For a sorted list, `add(newEntry)` must traverse the chain to locate the point of insertion. If the insertion is at the end of the chain, the traversal will give you a reference to the last node. A separate tail reference is not needed.

# Efficiency of the Linked Implementation

ADT Sorted List Operation	Array	Linked
<code>add(newEntry)</code>	$O(n)$	$O(n)$
<code>remove(anEntry)</code>	$O(n)$	$O(n)$
<code>getPosition(anEntry)</code>	$O(n)$	$O(n)$
<code>getEntry(givenPosition)</code>	$O(1)$	$O(n)$
<code>contains(anEntry)</code>	$O(n)$	$O(n)$
<code>remove(givenPosition)</code>	$O(n)$	$O(n)$
<code>toArray()</code>	$O(n)$	$O(n)$
<code>clear()</code> , <code>getLength()</code> , <code>isEmpty()</code>	$O(1)$	$O(1)$

FIGURE 16-5 The worst-case efficiencies of the operations on the ADT sorted list for two implementations



# Implementation That Uses the ADT List

- Sorted list a natural application for ADT list
- Possible ways
  - Use list as data field within class that implements sorted list
  - Use inheritance to derive sorted list from list
- Our class **SortedList** will implement the interface **SortedListInterface**
- View source code, [Listing 16-A](#)

An instance of a list

An instance of a sorted list

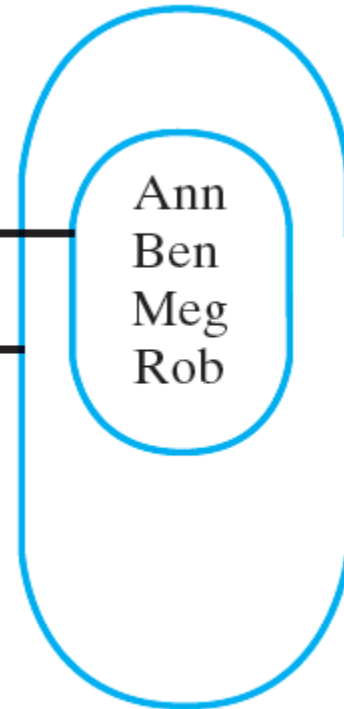


Figure 16-6 An instance of a sorted list that contains a list of its entries

Question 11 If a sorted list contains five duplicate objects and you use the below method `remove` to remove one of them, what will be removed from the list: the first occurrence of the object, the last occurrence of the object, or all occurrences of the object?

```
public boolean remove(T anEntry)
{
    boolean result = false;
    int position = getPosition(anEntry);
    if (position > 0)
    {
        list.remove(position);
        result = true;
    }
    return result;
}
```

Question 11 If a sorted list contains five duplicate objects and you use the below method `remove` to remove one of them, what will be removed from the list: the first occurrence of the object, the last occurrence of the object, or all occurrences of the object?

```
public boolean remove(T anEntry)
{
    boolean result = false;
    int position = getPosition(anEntry);
    if (position > 0)
    {
        list.remove(position);
        result = true;
    }
    return result;
}
```

The first occurrence of the object. Note that `getPosition` returns the position of the first occurrence of the entry within the list.

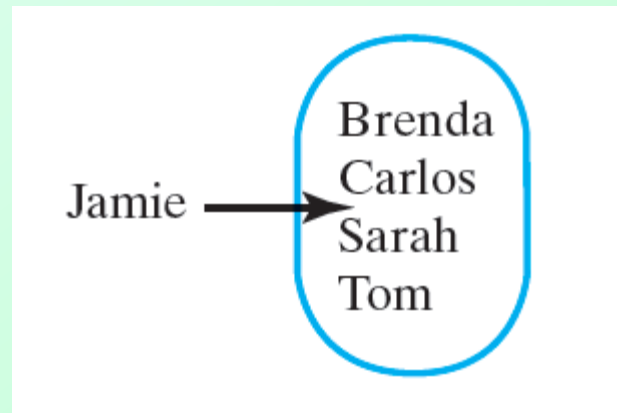


Figure 16-7 A sorted list in which *Jamie* belongs after *Carlos* but before *Sarah*

# Efficiency Issues

ADT List Operation	Array	Linked
getEntry(givenPosition)	$O(1)$	$O(n)$
add(newPosition, newEntry)	$O(n)$	$O(n)$
remove(givenPosition)	$O(n)$	$O(n)$
contains(anEntry)	$O(n)$	$O(n)$
toArray()	$O(n)$	$O(n)$
clear(), getLength(), isEmpty()	$O(1)$	$O(1)$

Figure 16-8 The worst-case efficiencies of selected ADT list operations for array-based and linked implementations

Question 15 Give an advantage and a disadvantage of using composition in the implementation of the class SortedList.

Question 15 Give an advantage and a disadvantage of using composition in the implementation of the class SortedList.

Advantage: The implementation is easy to write.

Disadvantage: The implementation is not efficient when the implementation of the underlying list is linked.



ADT Sorted List Operation	List Implementation	
	Array	Linked
add(new Entry)	$O(n)$	$O(n^2)$
remove(anEntry)	$O(n)$	$O(n^2)$
getPosition(anEntry)	$O(n)$	$O(n^2)$
getEntry(givenPosition)	$O(1)$	$O(n)$
contains(anEntry)	$O(n)$	$O(n)$
remove(givenPosition)	$O(n)$	$O(n)$
toArray()	$O(n)$	$O(n)$
clear(), getLength(), isEmpty()	$O(1)$	$O(1)$

Figure 16-9 The worst-case efficiencies of the ADT sorted list operations when implemented using an instance of the ADT list

# End

## Chapter 16

